

# Optimal usage of SSDs under Linux:

## Optimize your I/O Subsystem

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**Thomas-Krenn.AG**<sup>®</sup>

The server experts



# Introduction

## who I am

Werner Fischer



from Australia



Linux user  
since 2001



working for a  
Server vendor



freelancer  
IT journalist



Piano learner

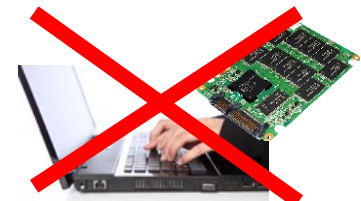


## who I am not

Kernel  
developer

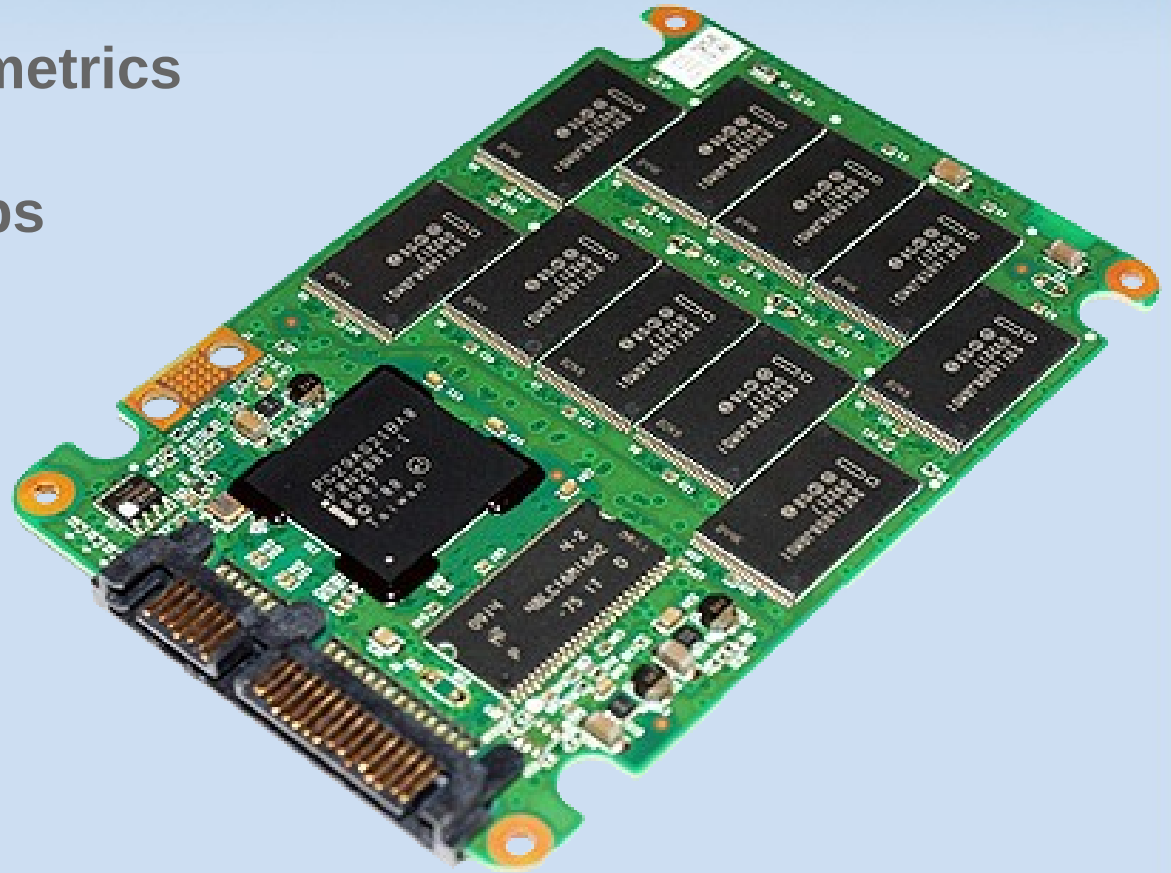


SSD  
developer



# Agenda

- 1) SSD layout
- 2) I/O performance metrics
- 3) Configurations tips

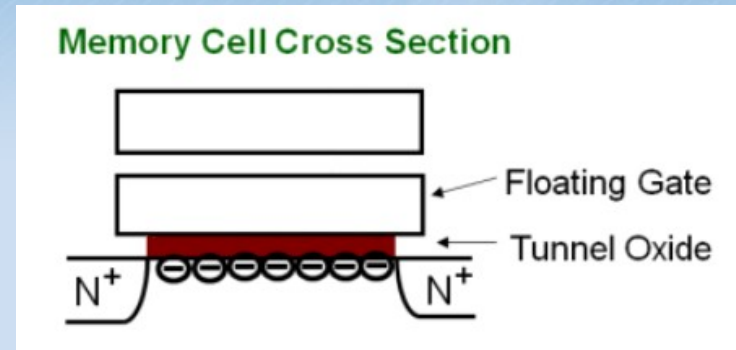


Source: maximumpc.com

# 1) SSD layout: memory cell

- **memory cells**

- NAND memory cell = MOS transistor with floating gate
- permanently store charge
  - programming puts electrons on floating gate
  - erase takes them off
- one program/erase (p/e) cycle is a round trip by the electrons
- back-and-forth round trips gradually damage the tunnel oxide
- endurance is limited, measured in number of p/e cycles:
  - 50nm MLC ~ 10.000 p/e cycles
  - 34nm/25nm/20nm MLC ~ 3.000 – 5.000 p/e cycles

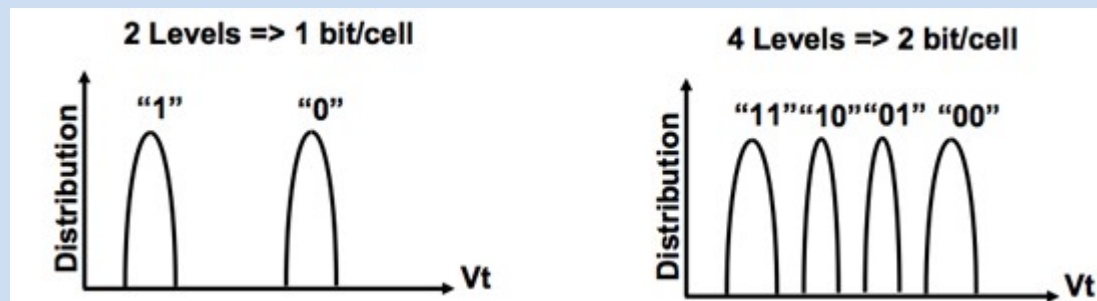


Source: Intel

# 1) SSD layout: memory cell

- **memory cells**

- SLC (Single Level Cell) → 1 bit per memory cell
- MLC (Multi Level Cell) → 2 bits per memory cell



Source: anandtech.com

- TLC (Triple Level Cell) → 3 bits per memory cell
  - 16LC (16 Level Cell) → 4 bits per memory cell
- **multiple memory cells (e.g. 16.384) build up a “page”**
    - page = smallest area, which can be read/written

# 1) SSD layout: page



- **one line = page within a SSD**
  - 8.192 Bytes (8 kiB)
  - can be read/written individually
  - cannot be changed/erased individually

Note: example sizes of pages and blocks are taken from Intel's Series 320 SSDs (with IMFT's 25nm Flash chips)

# 1) SSD layout: block



- **one line = page within a SSD**
  - 8.192 Bytes (8 kiB)
  - can be read/written individually
  - cannot be changed/erased individually
- **one blackboard = block within a SSD**
  - consists of 256 lines (pages), 2.097.152 Bytes (2 MiB)
  - smallest area which can be individually erased (we have only watering-cans for that ;-)



Note: example sizes of pages and blocks are taken from Intel's Series 320 SSDs (with IMFT's 25nm Flash chips)

# 1) SSD layout: change page



- **easier way to change lines?**



# 1) SSD layout: change page



- **easier way to change lines!**
  - (1) mark old line as invalid

# 1) SSD layout: change page



- **easier way to change lines!**
  - (1) mark old line as invalid
  - (2) store new content in a free line

# 1) SSD layout: change page



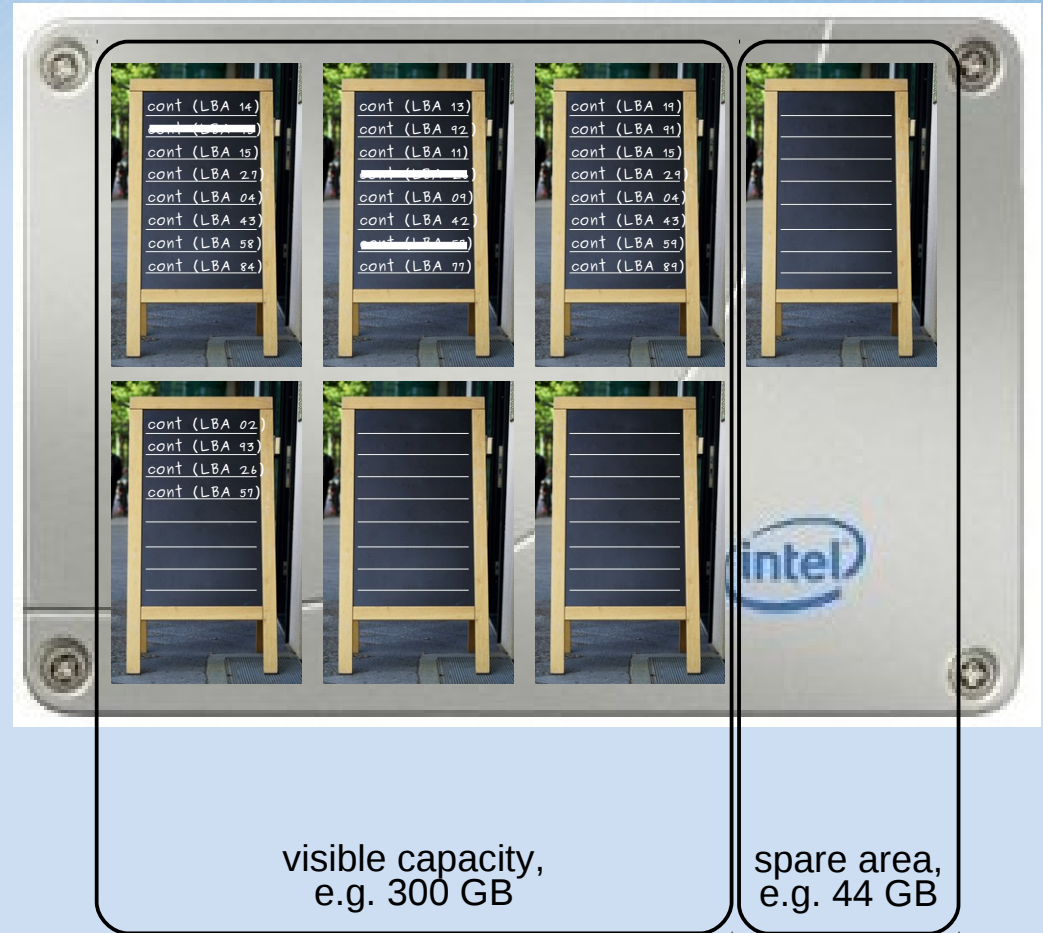
- **easier way to change lines!**
  - (1) mark old line as invalid
  - (2) store new content in a free line
- **this can be done as long as there are enough free lines left...**



we need spare blackboards to have enough free lines!

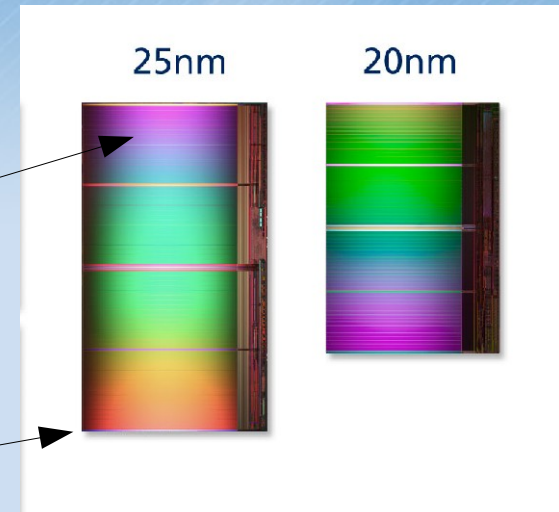
# 1) SSD layout: spare area

- **SSDs need spare area**
  - avoids the erasement of a block when a single page is changed
  - after some time spare area will be filled up, too
  - cleaning gets necessary (garbage collection)

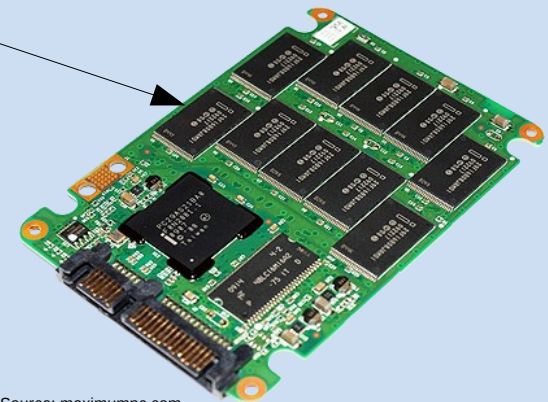


# 1) SSD layout: blocks → planes → dies → TSOPs

- **planes**
  - multiple blocks make up a plane
  - e.g. 1.024 blocks = 1 plane
- **dies**
  - multiple planes make up a die
  - e.g. 4 planes = 1 die
- **TSOP (thin small outline package)**
  - multiple dies (e.g. 1 - 8)
- **SSDs**
  - e.g. 10 TSOPs



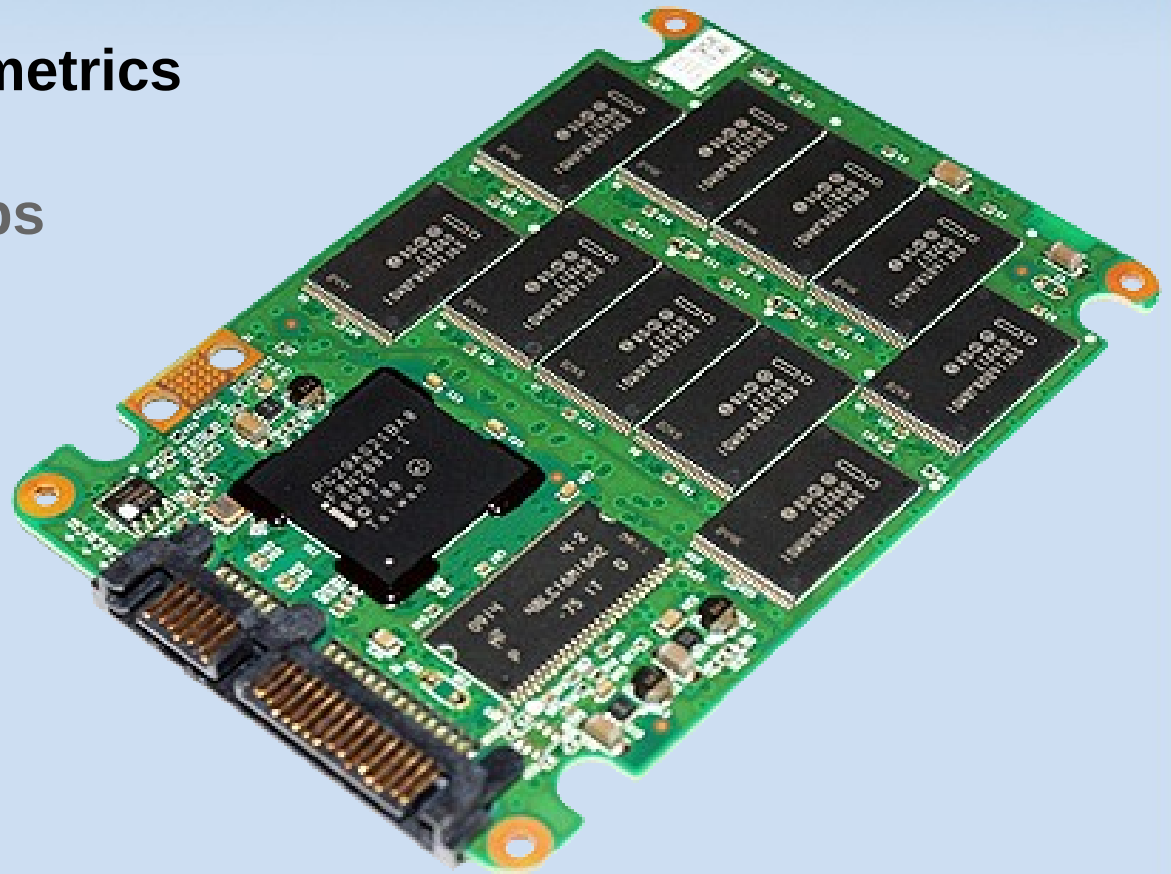
Source: [http://newsroom.intel.com/community/intel\\_newsroom/blog/2011/04/14](http://newsroom.intel.com/community/intel_newsroom/blog/2011/04/14)



Source: [maximumpc.com](http://maximumpc.com)

# Agenda

- 1) SSD layout
- 2) I/O performance metrics
- 3) Configurations tips



Source: maximumpc.com

## 2) I/O performance metrics

- **throughput**
  - MByte/s
  - throughput analogy:
    - # of persons/h from Berlin → Prague
- **# of I/O operations per second**
  - IOPS analogy:
    - # of individual trips to Prague (from Berlin, Vienna, Paris, Rome, ...)
- **latency because of queue depth**
  - queue depth analogy:
    - vehicles must use a ferry to reach destination
    - with how many vehicles does the ferry depart?



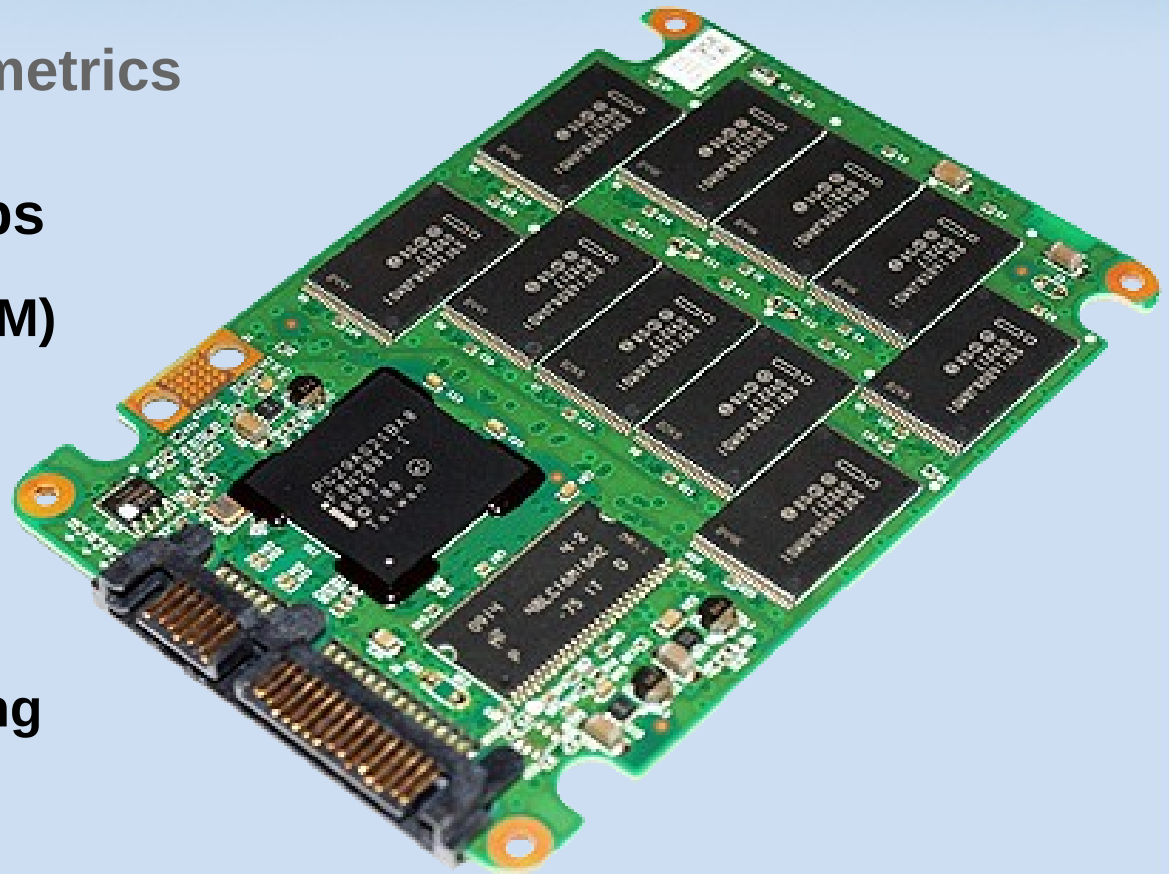
# Agenda

## 1) SSD layout

## 2) I/O performance metrics

## 3) Configurations tips

- AHCI (NCQ+DIPM)
- TRIM (discard)
- noatime
- tmpfs
- alignment
- over-provisioning



Source: maximumpc.com



### 3) Configuration tips: AHCI (NCQ+DIPM)

- **NCQ (Native Command Queuing)**
  - allows SSD to execute multiple I/O requests in parallel
  - boosts throughput
  - configure queue depth to get your optimal balance between max. # of IOPS and lowest latency

```
root@werner-t410: ~
root@werner-t410:~# hdparm -I /dev/sdb | grep -i queue
    Queue depth: 32
    *      Native Command Queueing (NCQ)
root@werner-t410:~# cat /sys/block/sdb/device/queue_depth
31
root@werner-t410:~# echo 5 > /sys/block/sdb/device/queue_depth
root@werner-t410:~# cat /sys/block/sdb/device/queue_depth
5
root@werner-t410:~#
```

### 3) Configuration tips: AHCI (NCQ+DIPM)

- **DIPM (Device Initiated Interface Power Management)**

- reduces idle power down to 0,1 Watt

```
root@werner-t410: ~
root@werner-t410:~# hdparm -I /dev/sda | grep -e "Device-initiated\|Enabled"
      Enabled Supported:
      *   Device-initiated interface power management
root@werner-t410:~#
```

- **Aggressive Link Power Management**

- min\_power
- medium\_power
- max\_performance

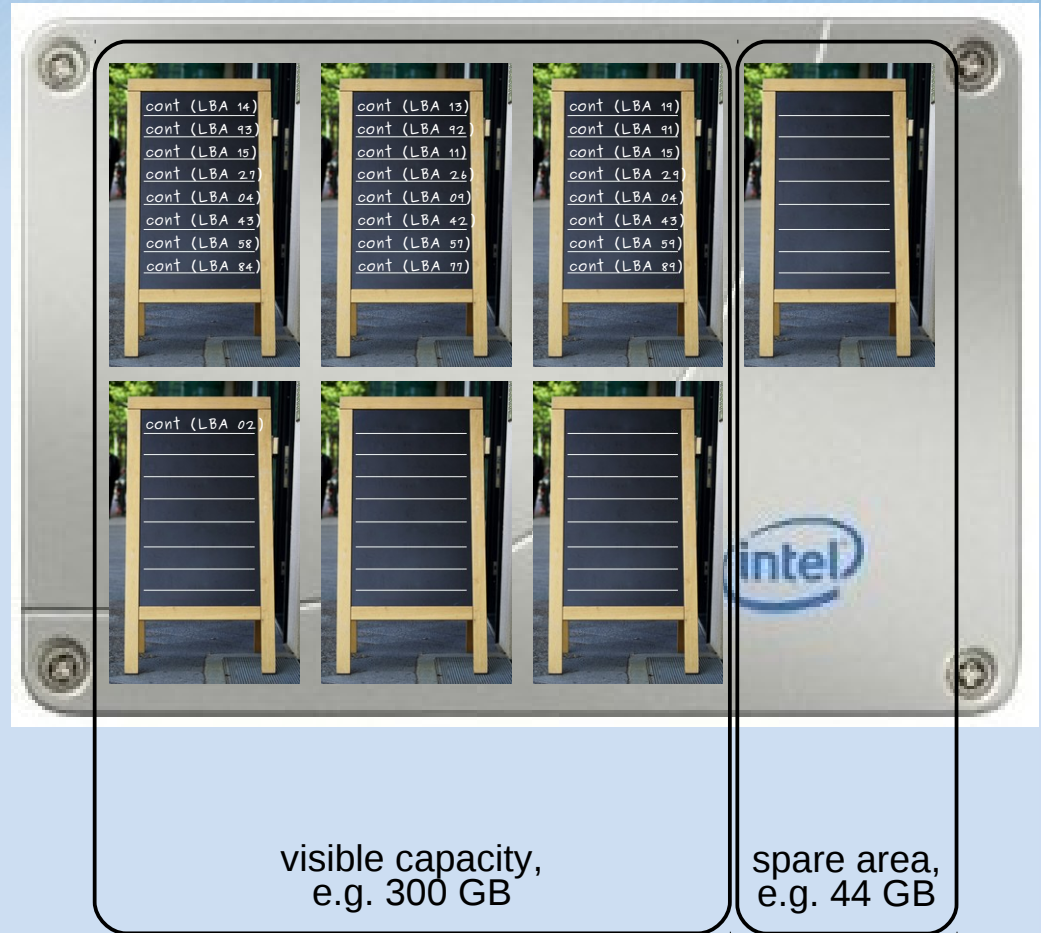
```
root@werner-t410: ~
root@werner-t410:~# lsscsi --classic
Attached devices:
Host: scsi0 Channel: 00 Target: 00 Lun: 00
  Vendor: ATA      Model: INTEL SSDSA2M160 Rev: 2CV1
  Type:   Direct-Access      ANSI SCSI revision: 05
Host: scsi1 Channel: 00 Target: 00 Lun: 00
  Vendor: HL-DT-ST Model: DVDROM GU10N   Rev: MX05
  Type:   CD-ROM             ANSI SCSI revision: 05
root@werner-t410:~# cat /sys/class/scsi_host/host0/link_power_management_policy
max_performance
root@werner-t410:~#
```

More Details: [http://docs.redhat.com/docs/en-US/Red\\_Hat\\_Enterprise\\_Linux/6/html/Power\\_Management\\_Guide/ALPM.html](http://docs.redhat.com/docs/en-US/Red_Hat_Enterprise_Linux/6/html/Power_Management_Guide/ALPM.html)

updated slide! (has been corrected after the talk)

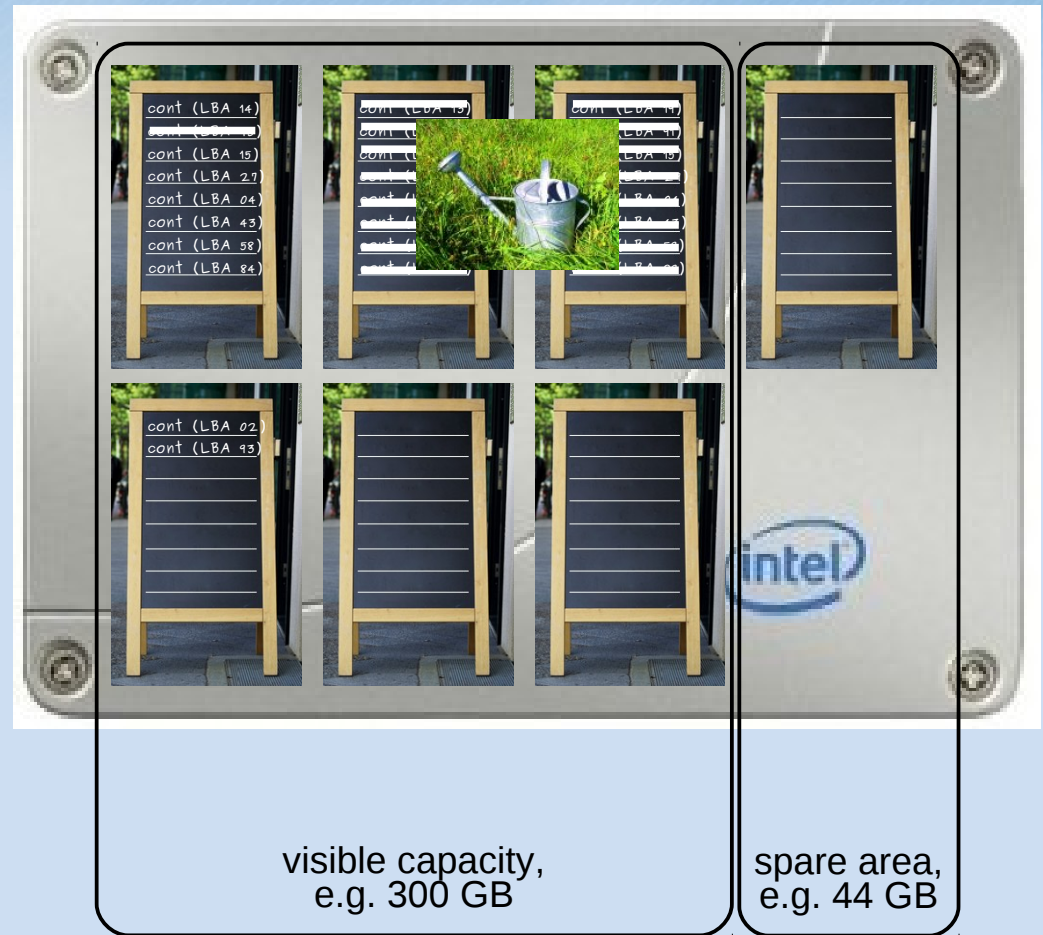
### 3) Configuration tips: TRIM (discard)

- **ATA TRIM**
  - tells SSD which data can be discarded



### 3) Configuration tips: TRIM (discard)

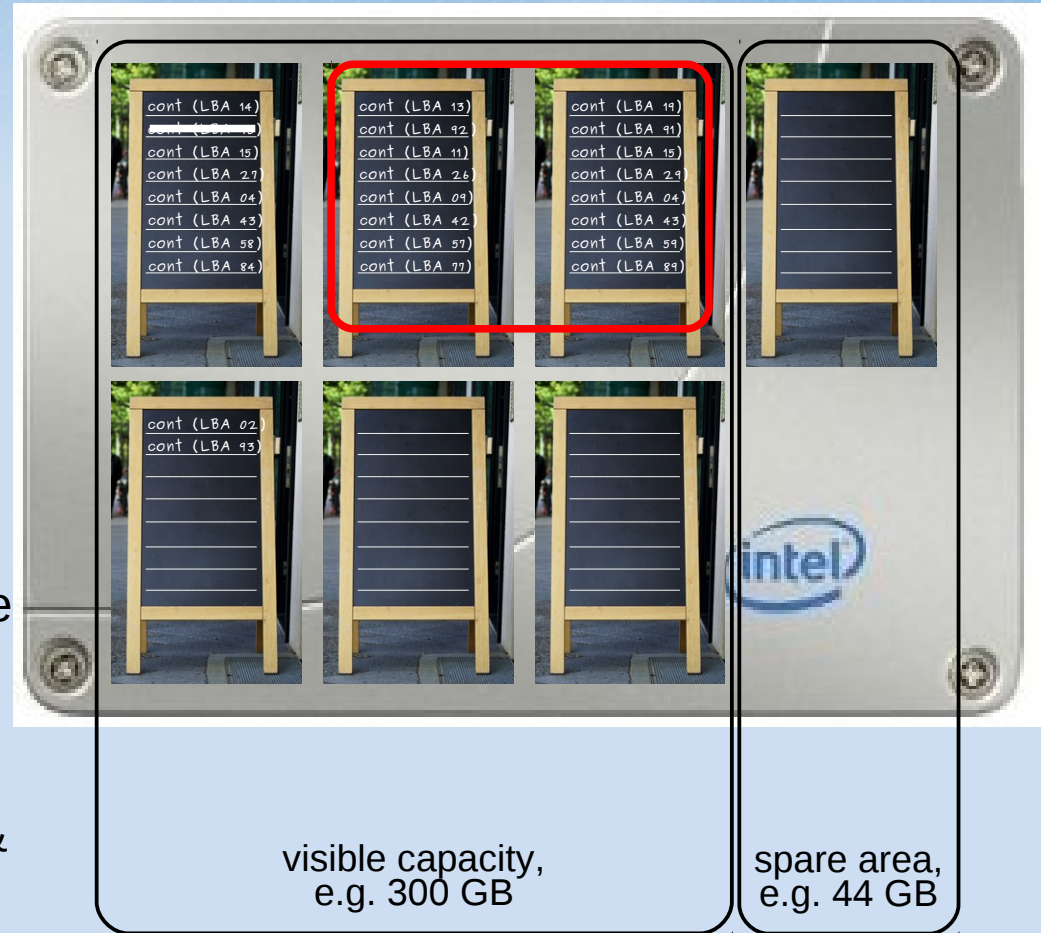
- **ATA TRIM**
  - tells SSD which data can be discarded



### 3) Configuration tips: TRIM (discard)

- **ATA TRIM**

- tells SSD which data can be discarded
- without TRIM:
  - deleting a big file (e.g. 100 GB) would lead to keep unusable data
  - unusable data will be maintained during garbage collection!
  - more overhead → lower performance & lower endurance!



## 3) Configuration tips: TRIM (discard)

- **ATA TRIM using discard infrastructure in Linux**
  - online discard
    - Ext4: since Kernel 2.6.33
    - XFS: since Kernel 3.0
    - Btrfs: since Kernel 2.6.32
  - batched discard (using fstrim command)
    - Ext4: since Kernel 2.6.37
    - XFS: since Kernel 2.6.38
    - Btrfs: since Kernel 2.6.39
  - pre-discard on format
    - E2fsprogs  $\geq$  1.41.10
    - Xfsprogs  $\geq$  3.1.0

### 3) Configuration tips: TRIM (discard)

- **ATA TRIM using discard infrastructure in Linux**
  - I/O stack discard support (device mapper):
    - since Kernel 2.6.36: DM targets delay, linear, mpath, stripe
    - since Kernel 2.6.38: DM mirror target
  - no I/O stack discard support yet:
    - MD raid
- **alternatives to discard: wiper.sh / raid1ext4trim.sh**
  - use `hdparm --trim-sector-ranges-stdin`
  - read warnings in the source of those scripts
  - cannot be used with device mapper

### 3) Configuration tips: TRIM (discard)

- “does TRIM work?” howto

```
root@werner-t410: ~
root@werner-t410:~# sudo hdparm -I /dev/sda | grep -i trim
*      Data Set Management TRIM supported (limit 8 blocks)
*      Deterministic read ZEROs after TRIM
root@werner-t410:~# echo "ABCD" > testfile; sync
root@werner-t410:~# hdparm --fibmap testfile

testfile:
  filesystem blocksize 4096, begins at LBA 61052928; assuming 512 byte sectors.
  byte_offset  begin_LBA      end_LBA      sectors
                0   73923472   73923479      8
root@werner-t410:~# hdparm --read-sector 73923472 /dev/sda | head -n 4

/dev/sda:
reading sector 73923472: succeeded
4241 4443 000a 0000 0000 0000 0000 0000
root@werner-t410:~# rm -f testfile
root@werner-t410:~# fstrim -v . ; sync
.: 3496751104 bytes was trimmed
root@werner-t410:~# hdparm --read-sector 73923472 /dev/sda | head -n 4

/dev/sda:
reading sector 73923472: succeeded
0000 0000 0000 0000 0000 0000 0000 0000
root@werner-t410:~#
```



### 3) Configuration tips: noatime & tmpfs

- **noatime / relatime**

- omits writes of metadata on every read

```
root@werner-t410: ~  
root@werner-t410:~# vi /etc/fstab  
root@werner-t410:~# sleep 10; date; grep noatime /etc/fstab  
Son Okt 23 20:53:50 CEST 2011  
UUID=1b7be627-fcb5-43b4-a9fb-72c112041c53 /          ext4    errors=remount-ro,noatime 0 1  
UUID=3bb73b03-8076-4d67-954c-6fee0c67080c /home      ext4    defaults,noatime          0 2  
root@werner-t410:~# stat /etc/fstab  
  File: `/etc/fstab'  
  Size: 965          Blocks: 8          IO Block: 4096   regular file  
Device: 803h/2051d  Inode: 394267      Links: 1  
Access: (0644/-rw-r--r--)  Uid: (   0/   root)  Gid: (   0/   root)  
Access: 2011-10-23 20:53:34.216717457 +0200  
Modify: 2011-10-23 20:53:34.216717457 +0200  
Change: 2011-10-23 20:53:34.216717457 +0200  
root@werner-t410:~#
```

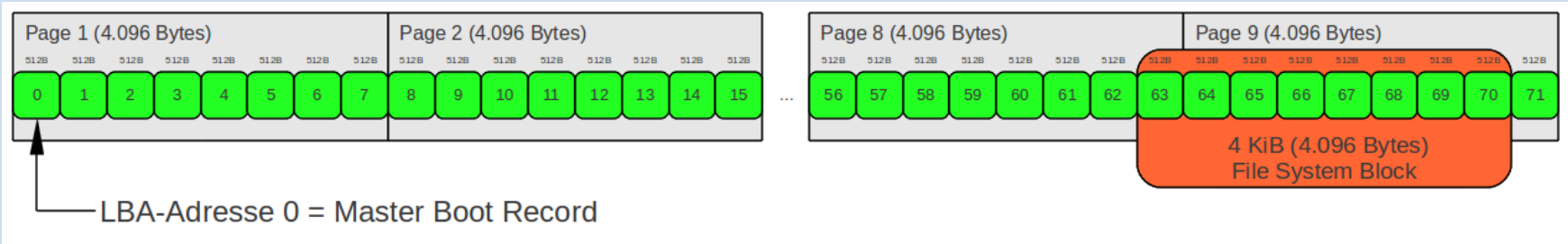
- **tmpfs**

- for temporary data like /tmp/, /var/tmp/, /var/cache/, ...

# 4) Configuration tips: alignment

- **align partition and file systems**

- wrong alignment:



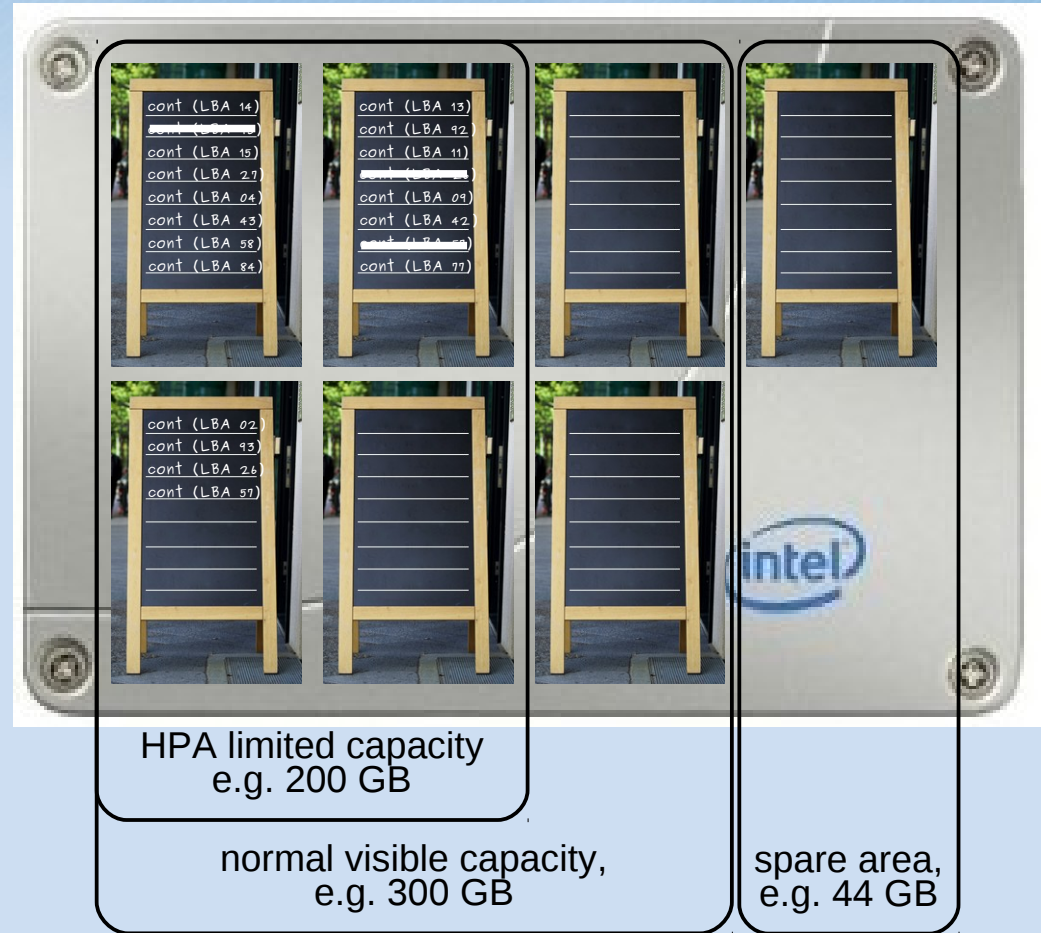
- use fdisk parameters: `fdisk -c -u /dev/sda`

- correct alignment:



## 4) Configuration tips: over-provisioning

- **do not use full normal visible capacity**
  - activate HPA (host protected area)
    - ATA8-ACS SET MAX ADDRESS
    - use `hdparm -N`
  - or simply do not partition full visible capacity
  - in either case if SSD has been used before:
    - do a secure erase to TRIM all blocks



## 4) Configuration tips: over-provisioning

- **over-provisioning is useful when discard cannot be used yet (e.g. MD-RAID, hardware RAID, ...)**
- **measurements by Intel:**



Source: Intel

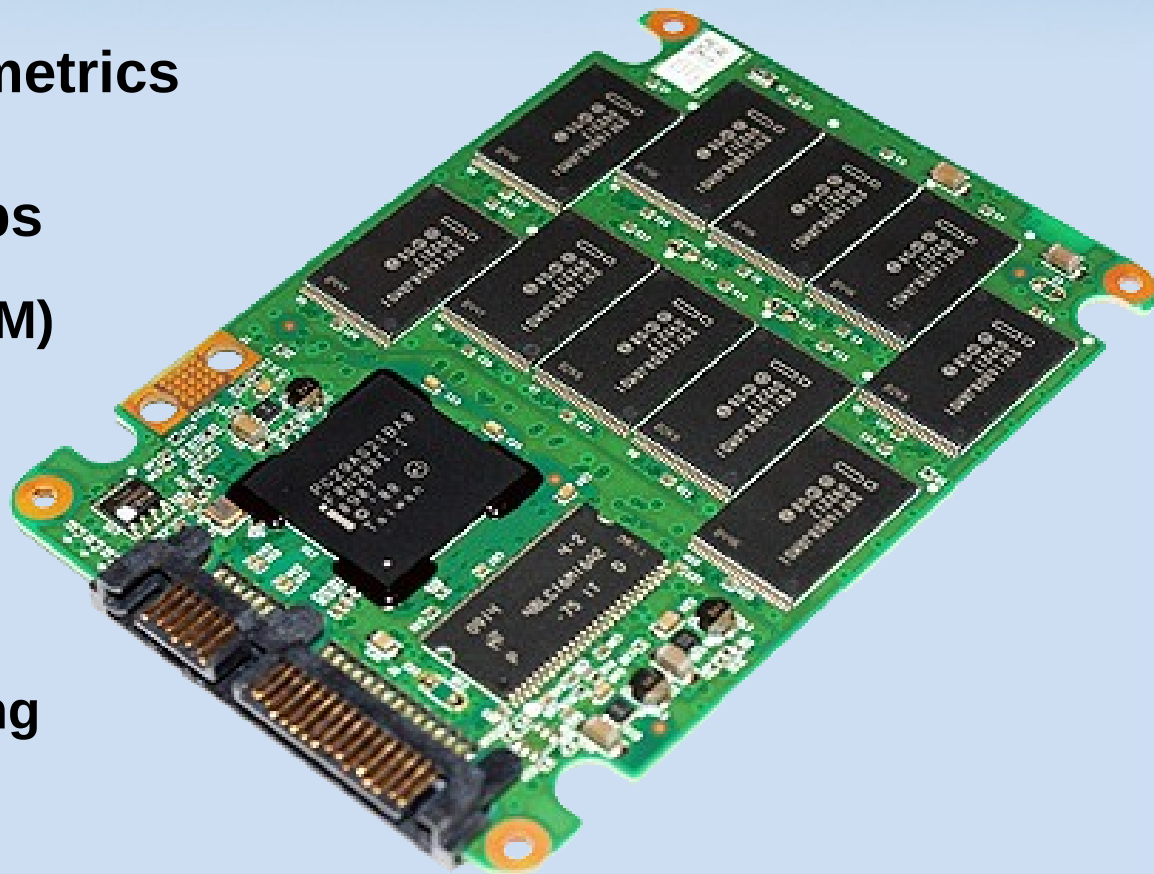
# Review:

## 1) SSD layout

## 2) I/O performance metrics

## 3) Configurations tips

- AHCI (NCQ+DIPM)
- TRIM (discard)
- noatime
- tmpfs
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# Thanks for your time!

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